Abstract—We discuss the capability of a new feature set for malware detection based on potential component leaks (PCLs). PCLs are defined as sensitive data-flows that involve Android inter-component communications. We show that PCLs are common in Android apps and that malicious applications indeed manipulate significantly more PCLs than benign apps. Then, we evaluate a machine learning-based approach relying on PCLs. Experimental validations show high performance for identifying malware, demonstrating that PCLs can be used for discriminating malicious apps from benign apps.

I. INTRODUCTION

Recent statistics from McAfee states that the total number of mobile malware samples exceeded five million, growing by 112% in one year. Indeed, mobile devices are a popular target for attackers, and app markets are still abused by malware developers for spreading their malicious apps. Consequently, the security guard of such markets has become an essential challenge for both end users and market maintainers.

Machine learning techniques, by allowing to sift through large sets of apps to detect malicious apps, appear to be promising for large-scale malware detection and eventually to keep malicious apps from entering app markets [2]. State-of-the-art machine learning approaches for Android malware detection mainly differ in the feature sets that are considered for training the classifiers. For example, Canfora et al. [7] rely on system calls and permissions while Gascon et al. [12] use call graph properties. Other examples of recurrent feature sets include Java code properties, Intent Filter information, strings, etc.

Recently, MUDFLOW [6] has proposed to extract behavioral features by taking into account sensitive data flows in Android apps to identify malware. In this paper, we also study the capability of specific sensitive data-flow features to be discriminative in Android malware detection as in MUDFLOW. Contrary to MUDFLOW, for which the source and sink of the data-flow are necessarily within a single component, we consider data-flows that may lead to leaks between two components such as data-flow coming from a source and going out of the component without knowing yet if the related data will go to a sink. Indeed, these potential component leaks (PCLs) are meaningful characteristics of malware since researchers have shown that the inter-component communication (ICC) mechanism introduces a lot of vulnerabilities (e.g., Activity Hijacking [21]).

In our previous work, we have developed PCLeaks [17], a tool for detecting potential component leaks involving two components. For the purpose of this study, we have extended PCLeaks to take into account the case where more than two components are involved in the leak (e.g., one component is used as a bridge component between two others).

This paper reports on an empirical investigation into potential component leaks in Android apps. Eventually, we assess the relevance of potential component leaks as features for machine learning-based Android malware detection. For instance, we experimentally check whether the performance achieved with these features can be generalized to different clusters of Android apps.

The contributions of this paper are as follows:

• We present a discussion on the different types of potential component leaks in Android apps.
• We empirically investigate the distribution of potential component leaks in malicious and benign app datasets.
• We further investigate the discriminative power of PCL-based features for machine learning-based malware detection.

Space limitations preclude us from including some experiments on generalization analysis of machine learning-based classifiers. We thus publish an accompanying Technical Report [15], which investigates the extent of dependency between the training data and the yielded classifiers. Through the investigation, we advocate that the assessment of feature sets must dig into the composition of training datasets.

II. POTENTIAL COMPONENT LEAKS

In a previous work [16], we have shown that ICCs are used to leak sensitive data across components. Any component can potentially participate in a leak, for instance by retrieving a piece of sensitive information, by sending this information, or simply by playing the role of a bridge between two other components. Thus, when performing static analysis on a single component, some of the data-flow paths, leaking data across the boundary of the component can be identified. In this paper, those data-flow paths are called Potential Component Leaks (PCLs). A PCL is not per se a leak but it might be exploited by other components and eventually contribute to a leak of private data.
Fig. 1 illustrates four different scenarios of data leakages, three of which represent PCLs. In this figure, (A) represents the “traditional” intra-component privacy leaks, which have been well studied in the literature [2]. In the present study, we do not consider such leaks since they are fully contained in one single component—i.e., a piece of data is both obtained and leaked inside one component—and hence do not involve any Inter-Component Communication. In (B), the leaked data is exfiltrated by the component, while in (C) the data is obtained from the component. Finally, in (D), the leaked data travels through the component which is merely used as a bridge. From these schematic examples, we see that a component is involved in a PCL either by providing an entry-point or by providing an exit-point for leaking the data.

(A)  
(B)  
(C)  
(D)  

Legend
- entry-point
- source
- sink
- exit-point

Fig. 1: Examples of leak schemes including “traditional” intra-component leaks (A) and PCLs ((B), (C) and (D)).

A. Characterization of Entry and Exit points

Unlike traditional Java apps, which come with a single entry point (main method), an Android app includes several components, each of which may contain several entry-points for launching the app. Because components can use different methods to call out other components, each component may contain multiple exit-points.

We now detail the criteria for identifying such entry-points and exit-points.

Entry-points: Entry-points are methods where data can be transferred into a component, through parameters, between components. In our study, we consider the following methods:
- Any method such as `getStringExtra()` that obtains data from Intents (or Bundles).
- Any lifecycle method that takes an Intent as a parameter.
- Any method that obtains data from `ContentValues`.
- Any method of `ContentResolver` such as `query()` that acquires data from other components (or apps).

Exit-points: An exit-point is a method call through which data can be transferred outside a component. For example, the `startActivity()` method can be used to trigger data exchange when one component launches another. We consider all such ICC methods as exit-points. We also take into account such methods of `ContentResolver` such as `insert()` that are capable to transfer data to `ContentProviders`.

2 It is not necessary for entry-points to get data from Intents since Intents can be directly leaked through components, e.g., type (D) in Fig. 1.

3 Like Intents, `ContentValues` are used to exchange data between components. However, they are used to transfer data to `ContentProviders` while Intents are used for the other three types of components.

B. PCL types

In this section, we introduce the three types of PCLs that are investigated. An example of PCL is shown in Fig. 2 where a sensitive data (device id) is collected in a first component (1) and leaked through an ICC method to another component (2) which simply forwards it to a third component (3) where it is eventually leaked outside the device through SMS. To detect such leaks, PCLeaks performs static taint analysis on each app and tracks the sensitive data across components from its source to sink.

Potential Active Component Leak (PACL). We define a PACL as a taint flow path starting from a source (defined as calls into resource methods returning non-constant values into the application code) and ending with an exit-point. Such PCLs are referred to as “active”, as the involved component is actively leaking sensitive data that it collected itself to other components (cf. (1) in Fig. 2).

Potential Bridge Component Leak (PBCL). We define a PBCL as a taint flow path starting from an entry-point and ending with an exit-point. Such PCLs are referred to as “bridge”, as the involved component is transferring sensitive data collected by a different component to another component (cf. (2) in Fig. 2).

Potential Passive Component Leak (PPCL). We define a PPCL as a taint flow path starting from an entry-point and ending with a sink (defined as calls into resource methods accepting at least one non-constant data value from the application code as parameter, if and only if a new value is written or an existing one is overwritten on the shared resource (e.g., GSM network) [23]). Such PCLs are referred to as “passive”, as the involved component is passively linking sensitive data collected by other components (cf. (3) in Fig. 2).

Note that the Android system provides two types of mechanism to protect components of being misused by other components: the export attribute and permissions. i) The export attribute is used to express the fact that other components can “access” the exported one. Thus, only exported components can potentially leak private data (PPCL and PBCL). ii) Permissions can be used at component level. When a component is protected by a permission, the apps that want to access this component must have first requested, and be granted this permission. In our detection of PCLs, we take into account these two types of mechanism, for instance by checking whether the export attribute is used or not.
III. Experimental Setup

In this section we detail the settings used in PCLeaks to yield PCLs. We also present the dataset for the experiments as well as the construction of the feature vectors for machine learning experiments.

A. PCLeaks Settings

In this study, all PCLs are detected by a new version of PCLeaks [17] which has been extended to detect PBCLs as well. PCLeaks uses a static taint analysis approach to track data paths from sources to sinks. We discuss in this section how such sources and sinks are determined in our work.

Sensitive sources and sinks. The key idea behinds static taint analysis is to identify a path that starts from a sensitive source and ends with a sensitive sink. In this study, the sensitive source and sink we use are extracted by SUSI [23], which automatically classifies all methods in the whole Android API as source, sink or neither. In Android 4.2, SUSI yields 18,076 source methods and 8,314 sink methods. Theoretically, there are 150,283,864 (18,076 * 8,314) different taint paths (the source to sink pairs) PCLeaks can report. Instead of identifying PCLs by method pairs, we group methods with similar functionality into categories (e.g., group methods Log.e() and Log.v() into category LOG) and use this category in lieu of the fully-qualified method name. This categorization allows to vastly reduce the number of different identifying pairs down to a more manageable value. We use the categories provided by SUSI for our study. Note that we have ignored methods that are classified as NO_CATEGORY by SUSI except for methods related to shared preferences since they are well used in Android apps, for which we create a new category (SHARED_PREFERENCES). Besides, we create four new categories (one for each component type) to further break down the behavior of potential component leaks.

Analysis Settings. As mentioned in [17], PCLeaks leverages the static taint analysis tool FlowDroid to identify data flows in Android apps. Because FlowDroid analyzes a whole Android app and aims to provide highly precise results, it usually takes a lot of time and resources to analyze an app. In favor of a faster analysis, we use the same FlowDroid settings as MUDFLOW [6] chooses (Explicit flow only, Disable flow-sensitive alias search, Maximum access of path length of 3, No-Layout mode and No static fields), which sacrifices some amount of precision for speed and memory. As a result, the detected potential component leaks may have false positives as well as false negatives. However, our goal in this paper is not to prove the presence or absence of flows but to study the distribution difference of potential component leaks between malware and goodware.

Advertisement Libraries. Most Android apps are free, they usually use advertisement to get profit, which are delivered through specific advertisement libraries. These libraries access sensitive data such as the unique device id to deliver personalized advertisements. However, the potential component leaks (flows) introduced by advertisement libraries are separate from the actual app code. As shown in MUDFLOW [6], advertisement libraries are frequently used and their flows (PCLs) thus become “normal”, diluting the impact of actual app flows. Therefore, we follow MUDFLOW’s assumption that advertisement libraries are trustworthy and ignore all the PCLs taking place in advertisement libraries, allowing our study to focus on the actual app PCLs. We use the same list of libraries MUDFLOW uses to exclude PCLs.

B. Datasets

For the purpose of our experiments, we collected a dataset of Android apps from Android markets including the official GooglePlay store. For each app, we also retrieved analysis results of anti-virus products hosted by VirusTotal. Then, based on the results of VirusTotal, we build two disjoint sets: One set, noted M (for Malware), containing only malicious apps and G (for Goodware) containing only benign apps. Each dataset contains 5,000 apps which we evaluate with PCLeaks.

All our experiments are performed on the UL HPC platform [24]. For each Android app, we allocate one core for PCLeaks to analyze it. The Java heap is set to 8 gigabytes and the time out is set to 12 hours. Recall that we start with two data sets containing 5,000 apps each for this study. Because some of them fail (e.g., exception or time out) or do not contain any PCLs, the result of the PCL extraction process contain 2,822 goodware and 3,785 malware, each containing at least one PCL.

C. Feature Set

One goal of this study is to assess if PCLs can be used as features for machine learning-based malware detection to suggest potential malicious apps. Machine learning algorithms cannot directly work on Android apps. Each app must be represented by a vector of properties, called a feature vector in the context of machine learning. In this study, our feature vectors are built with the results (PCLs) of PCLeaks, after analyzing all the apps in our dataset. Let L be the feature vectors we build, given an app a, for each PCL l_i ∈ L, we value it as either 0 for the case that a does not contain l_i or the actual number of l_i reported by PCLeaks. Recall that we use categorizations instead of methods to describe the detected taint flows (PCLs). Thus, our feature set is made of category pairs.

IV. Empirical Results

In this study, we address the following research questions:

- RQ1: Are PCLs common in Android apps?
- RQ2: Is there a significant difference in the presence of PCLs between malicious and benign apps? If so, is this difference similar for all PCL types?

4Explanations for these FlowDroid settings can be found at https://github.com/secure-software-engineering/soot-infoway-android/wiki
- RQ3: Can PCLs be used as features for machine learning-based malware detection?

A. Occurrence of PCLs in Android Apps

Fig. 3a plots the distribution of the number of PCLs per app from our dataset. The median value indicates that half of the apps contain at least 20 PCLs. Excluding outliers, which are automatically identified by the R statistics tool, the number of PCLs per app ranges from 0 to about 100. Because the various apps in our dataset are not equivalent in terms of code size and in terms of components, we further investigate the distribution of PCLs by normalizing the result in those two dimensions. Fig. 3b depicts the distribution of PCLs per 100 kilobyte of bytecode. The median number of PCLs per 100kB is around 3, while the maximum is slightly above 20. Finally, Fig. 3c presents the number of PCLs per component in the apps of the datasets. Components have a median value of 1 PCL, with a maximum of 6 PCLs per component.

B. Distribution of PCLs between malicious and benign apps

Following the findings on the occurrence of PCLs in Android apps in general, we further investigate whether the distribution of PCLs varies between malicious and benign apps. We therefore separately show in Fig. 4 the boxplots representing the number of PCLs for the malware and goodware datasets. The median values indicate that in general half of malware apps contain each more than 22 PCLs while this median value amounts to 8 for in goodware apps. To assess the significance of this difference we perform a Mann-Whitney-Wilcoxon (MWW) test which was successful (with p-value < 0.001).

We then explored whether this difference is similar for all types of PCLs considered in this study. The results show that PACLs, depicted in Figure 4b, are the most unequally distributed between malware and goodware. The median value for PACLs is 15 for malicious apps and only 1 for benign apps. The differences, according to MWW test, although statistically significant, are less important for PPCLs (Figure 4c, median values 4 for malware and 2 for goodware), and PBCLs (Figure 4d, median values 2 for malware and 0 for goodware).

C. Malware identification

Empirical findings from previous section on the presence of PCLs in malware and goodware datasets suggest that PCLs can be used to discriminate malicious apps from benign apps. In this section, we investigate this possibility by implementing and assessing a machine learning-based malware detection approach leveraging PCLs as classification features.

We perform extensive experiments, tuning different machine learning approach parameters, to gather insights for the practical use of PCLs as features. In particular we evaluate the performance of the features in combination with different machine learning classification algorithms. We also consider the impact of class imbalance in the dataset by varying the ratio between malware and goodware in the validation experiments.

Effect of Classification Algorithm. Fig. 5 plots the ROC graphs for the performance of the malware detector with different classification algorithms. All five algorithms yield an Area Under Curve (AUC) above 0.8, indicating good performance. The RandomForest algorithm achieves the best performance, although the overall performance of all algorithms are similar. This result suggests that the PCL-based feature set is not tailored for a specific algorithm.

Malware/Goodware Ratio. We investigate in detail how class imbalance in the constructed dataset threatens the performance of PCL-based malware classification. To this end, we customize three datasets, composed of 2,400, 3,600 and 4,800 apps with a malware/goodware ratio of 1, 2 and 3 respectively. We found that the performance decreases (but still with F-measure above 0.86) with the ratio of malware in the set. Such a finding was already shown in Allix et al.’s large scale empirical study with a different feature set [2].
In this paper, we have studied the distribution of potential component leaks, and used this information to detect malicious apps. This work is related to many existing techniques that leverage static taint analysis to detect privacy leaks, to detect malware or to perform empirical study on Android apps.

Information flow analysis. Information flow analysis has been well studied in Android community to detect vulnerabilities of Android apps. For instance, FlowDroid [5] performs “context-, flow-, field-, object-sensitive and lifecycle-aware static taint analysis for Android apps” to detect intra-component sensitive data flows. Several other works have been presented to detect inter-component information flows [13], [14], [16], [25]. For example, IccTA [16] leverages FlowDroid to perform inter-component static taint analysis through instrumenting Android apps, reducing an inter-component problem to an intra-component problems. Other techniques dynamically analyzes information flows of Android apps. For example, TaintDroid [10], one of the most sophisticated dynamic taint tracking system, uses a modified Dalvik virtual machine to track flows of private data.

In this work, we investigate potential component leaks, the component-based information flows, which are different from the above approaches. Our results are generated by our previous work, PCLeaks, which performs information flow analysis through the known ICC vulnerabilities (e.g., Activity Hijacking). CHEX [18] and ContentScope [27] are two other tools that tackle potential component leaks, however CHEX limits itself to only considering leaks related to Activity hijacking while ContentScope only takes into account leaks related to Content Provider.

Machine learning based malware detection. Recently, Avdiienko et al. presented an approach [6] closely related to ours. Both their approach and ours take sensitive data flows as features for machine learning-based malware detection, and both rely on FlowDroid to extract sensitive data flows. However, instead of taking into account all intra-component leaks, we focus on component-based privacy leaks, the so-called potential component leaks. Besides, we take into account SharedPreferences in our study, which has not been considered in Avdiienko et al. approach. Furthermore, we have investigated the clustering impact of the training data set, which at the moment is rarely investigated in the literature.
Allix et al. [2] empirically investigated the assessment of machine learning-based malware detectors for Android apps to measure the impact of datasets size and goodware/malware ratio, and the importance of validation scenarios. Our work is related in that we also measure the impact of several parameters and we raise one more factor to take into account when evaluating a malware detection approach: One specific approach may perform well only on a subset of Android applications.

Several other candidate features have been proposed to classify Android malware by using machine learning. For example, Peng et al. [22] apply probabilistic learning methods to the permissions of apps to detect malware. Gascon et al. [12] make use of embedded call graphs to build a malware detector. Other approaches [1, 4, 8, 11, 28] that rely on static or dynamic analysis also provide possible features for malware detection. Those features, along with the features we studied in this paper, could be combined to perform more accurate malware detection.

**Empirical study on Android apps.** In this work, we have empirically studied the distribution of potential component leaks. Empirical study provides a way of gaining knowledge quantitatively and qualitatively. Li et al. [16] present an empirical study on how Intent is used in Android apps, showing that Intent is commonly used in Android apps. Ruiz et al. [20] show the prevalence of multiple advertisement libraries in Android apps. Liu et al. [20] studied on the safety of storing non-shared data on public storage of Android. Egele et al. [9] illustrate that 10,327 out of the 11,748 apps they studied contain at least one mistake in their usage of cryptographic APIs. Maji et al. [19] perform fuzz testing to evaluate the robustness of Android ICC mechanism, showing that exception handling is rarely used and that it is possible to crash an app at runtime from an unprivileged user process. Allix et al. [3] perform a forensic analysis of Android apps, showing evidences that many Android malicious apps are developed at an industrial scale.

**VII. Conclusion**

In this study, we empirically investigated a new feature set for Android malware detection. This new feature set is based on potential component leaks (PCLs), which we define as sensitive data-flows that involve Android inter-component communications. We first showed that PCLs are common in Android apps. Then, further investigation showed that malicious apps contain significantly more PCLs than benign apps. Finally, we successfully applied PCLs as features for machine learning-based malware detection.

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